

# Ramsey theory of homogeneous structures

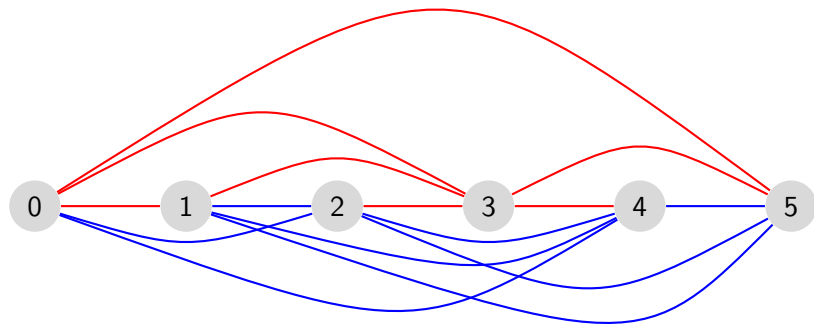
Natasha Dobrinen  
University of Denver

Notre Dame Logic Seminar

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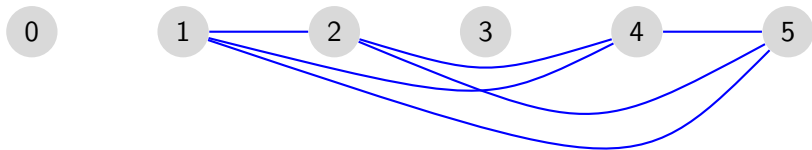
# Ramsey's Theorem for Pairs of Natural Numbers

Given a coloring of pairs of natural numbers into red and blue:



# Ramsey's Theorem for Pairs of Natural Numbers

There is an infinite subset  $M$  such that all pairs of numbers in  $M$  have the same color.



This can also be stated in terms of finding a complete infinite graph with all edges having the same color.

## Ramsey's Theorem and Logic

**Theorem.** (Ramsey, 1929) Given  $k, r \geq 1$  and a coloring

$$c : [\mathbb{N}]^k \rightarrow r,$$

there is an infinite  $M \subseteq \mathbb{N}$  such that  $c$  takes only one color on  $[M]^k$ .

This theorem appears in Ramsey's paper, *On a problem of formal logic*, and is motivated by Hilbert's *Entscheidungsproblem*:

Find a procedure for determining whether any given formula is valid.

Ramsey applied his theorem to solve this problem for formulas with only universal quantifiers in front ( $\Pi_1$ ).

One direction of extending Ramsey's Theorem is to trees.

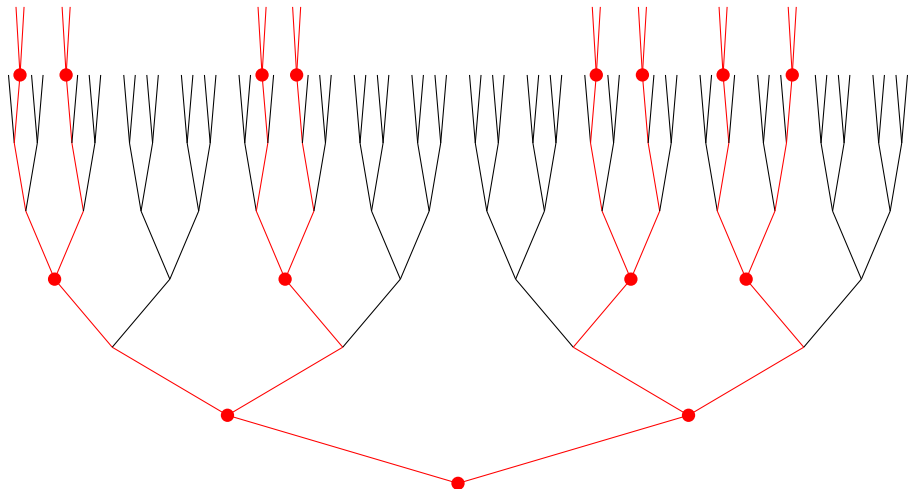
Several routes have been taken, and we will concentrate on the one on strong trees.

# Ramsey Theory for Strong Trees

A tree  $S \subseteq T = 2^{<\omega}$  is an **infinite strong subtree of  $T$**  iff it is (strongly) isomorphic to  $2^{<\omega}$ .

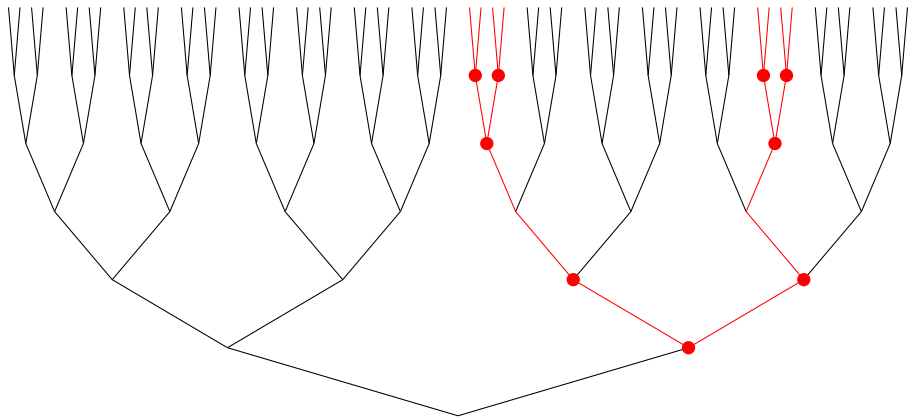
(This is a special case, but is sufficient for this talk.)

# Example, An Infinite Strong Subtree $S \subseteq 2^{<\omega}$



The nodes in  $S$  are of lengths  $0, 1, 3, 6, \dots$

## Example, An Infinite Strong Subtree $T \subseteq 2^{<\omega}$



The nodes in  $T$  are of lengths  $1, 2, 4, 5, \dots$



## Halpern-Läuchli Theorem - strong tree version for 2 trees

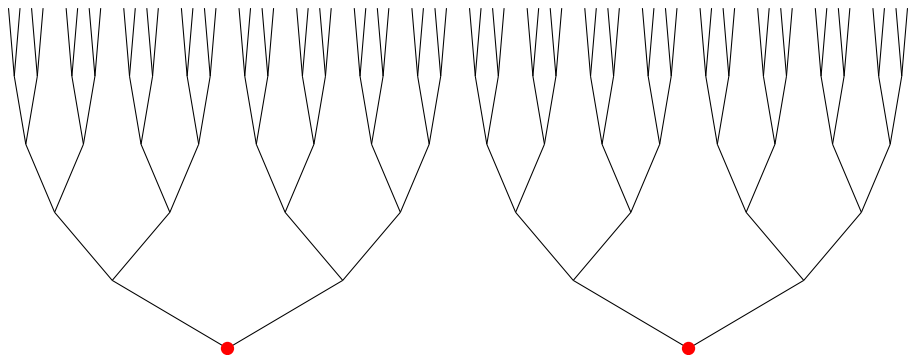
**Theorem.** (Halpern-Läuchli 1966) Let  $r \geq 2$ , and  $T_0 = T_1 = 2^{<\omega}$ . Given a coloring of the product of level sets of the  $T_i$  into  $r$  colors,

$$c : \bigcup_{n < \omega} T_0(n) \times T_1(n) \rightarrow r,$$

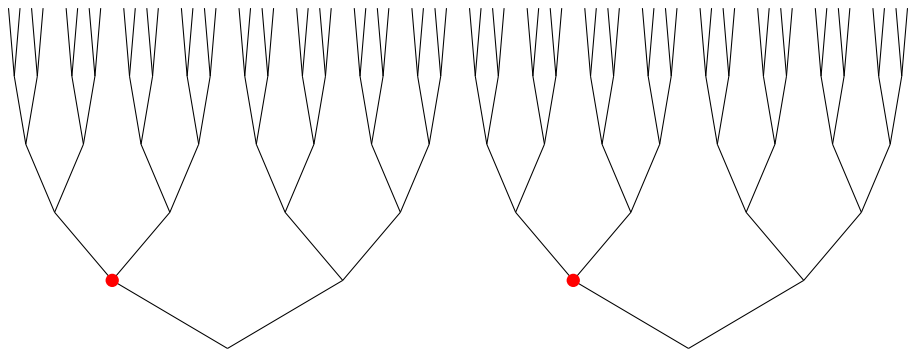
there are infinite strong trees  $S_i \leq T_i$  and an infinite sets of levels  $M \subseteq \omega$  where the splitting in  $S_i$  occurs, such that  $f$  is constant on  $\bigcup_{m \in M} S_0(m) \times S_1(m)$ .

This was found as a key lemma while proving that the Boolean Prime Ideal Theorem is strictly weaker than the Axiom of Choice over ZF. (See [Halpern-Lévy 1971].)

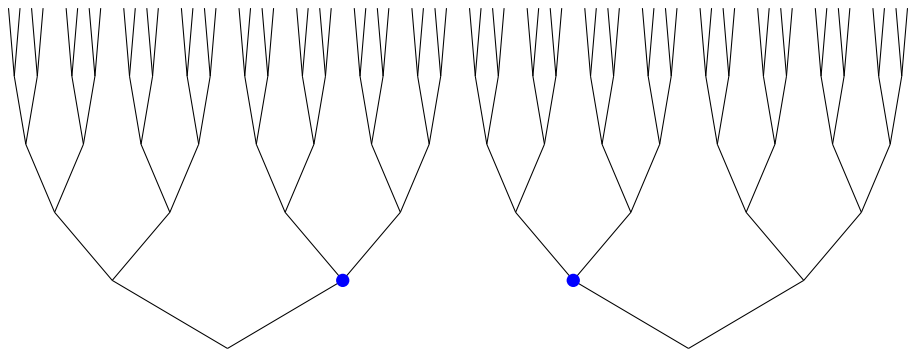
# Coloring Products of Level Sets: $T_0(0) \times T_1(0)$



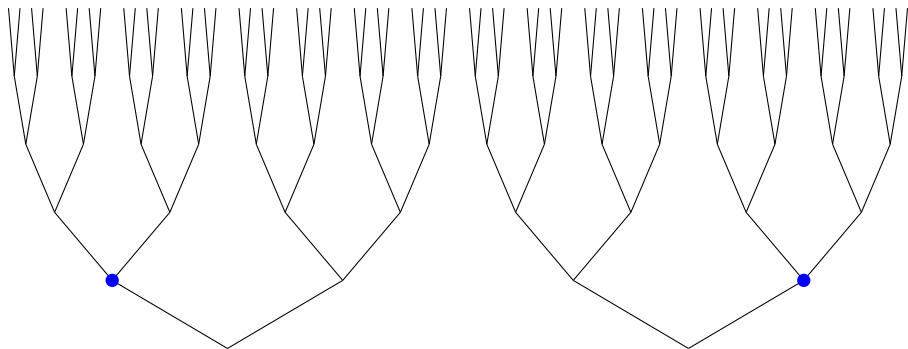
# Coloring Products of Level Sets: $T_0(1) \times T_1(1)$



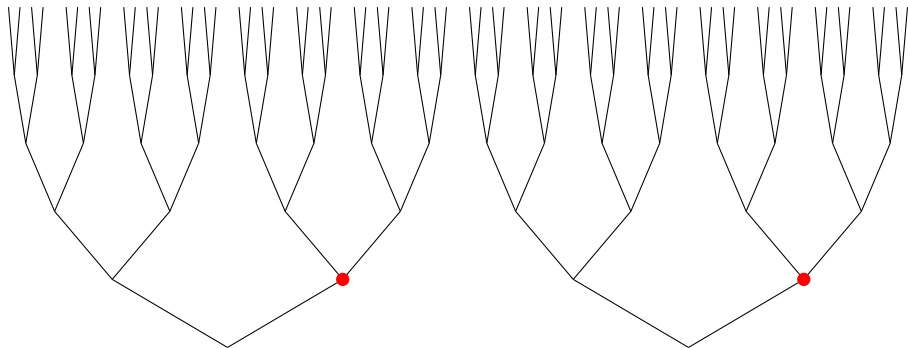
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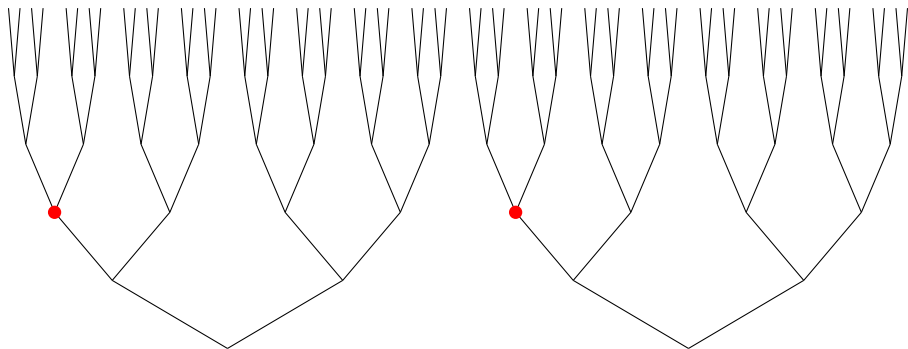
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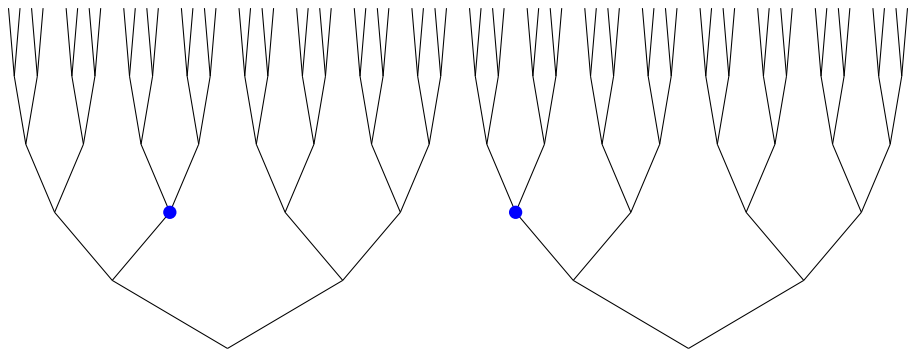
# Coloring Products of Level Sets: $T_0(1) \times T_1(1)$



## Coloring Products of Level Sets: $T_0(2) \times T_1(2)$

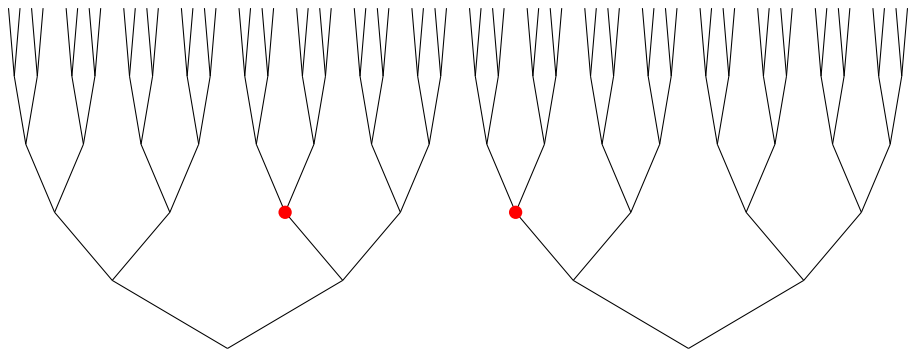


# Coloring Products of Level Sets: $T_0(2) \times T_1(2)$

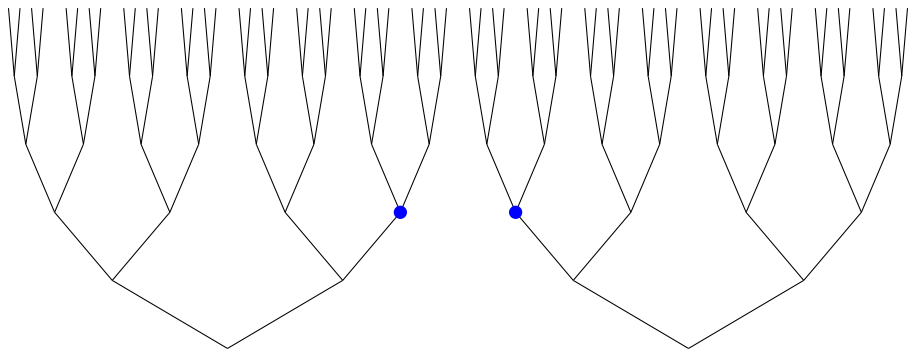




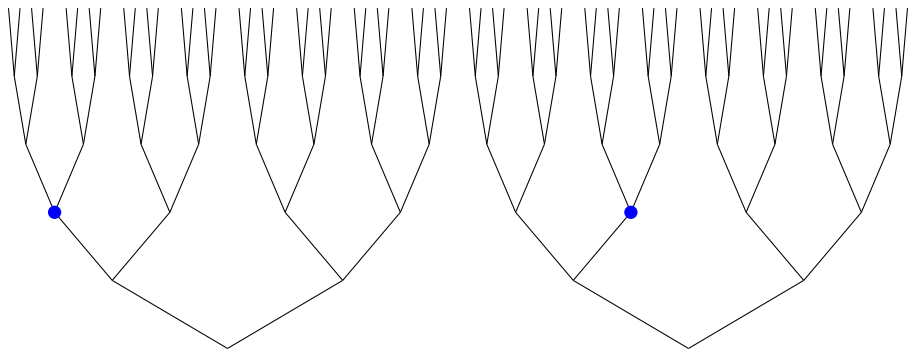
## Coloring Products of Level Sets: $T_0(2) \times T_1(2)$



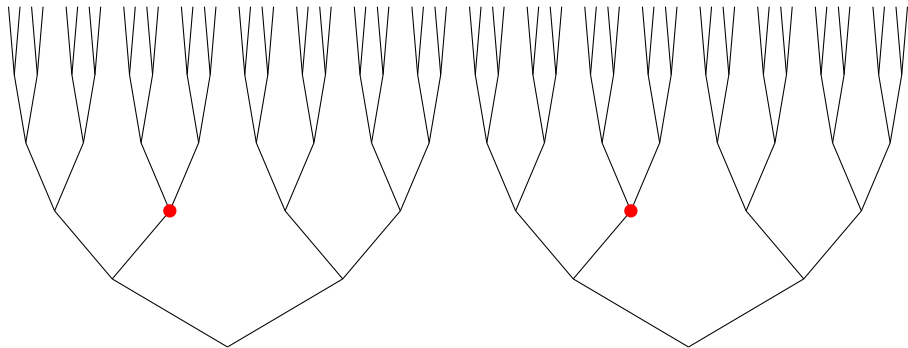
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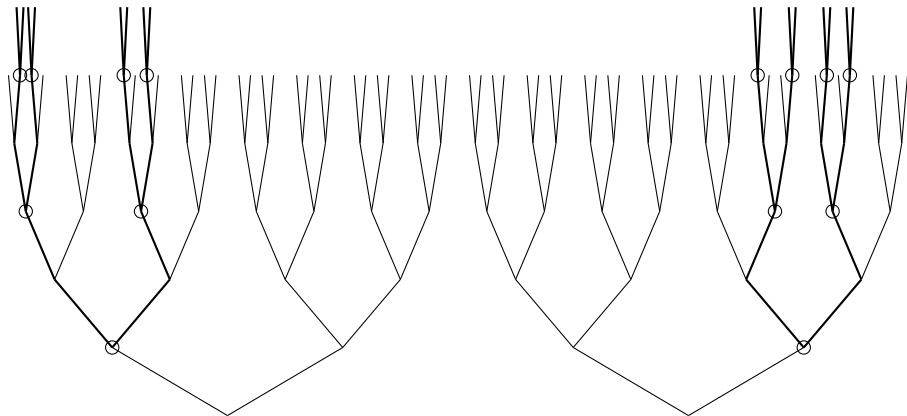


## Coloring Products of Level Sets: $T_0(2) \times T_1(2)$



Etc.

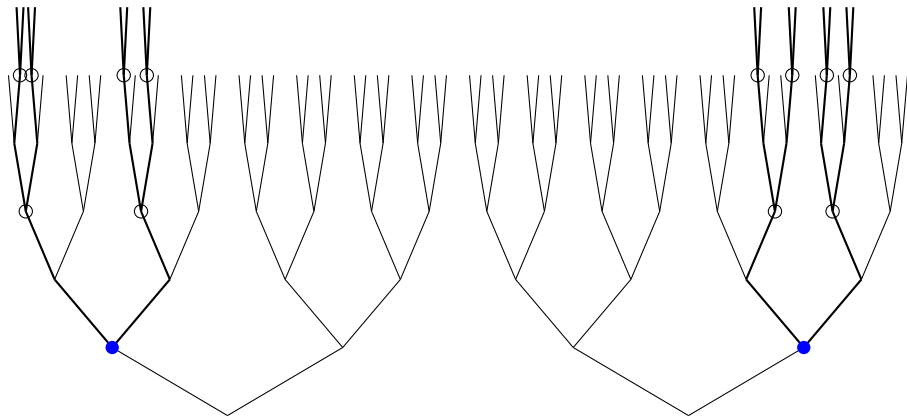
# HL gives Strong Subtrees with 1 color for level products



$S_0$

$S_1$

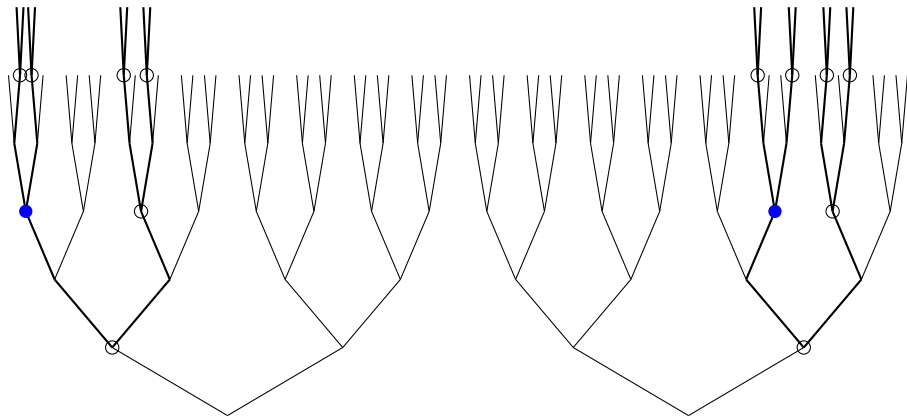
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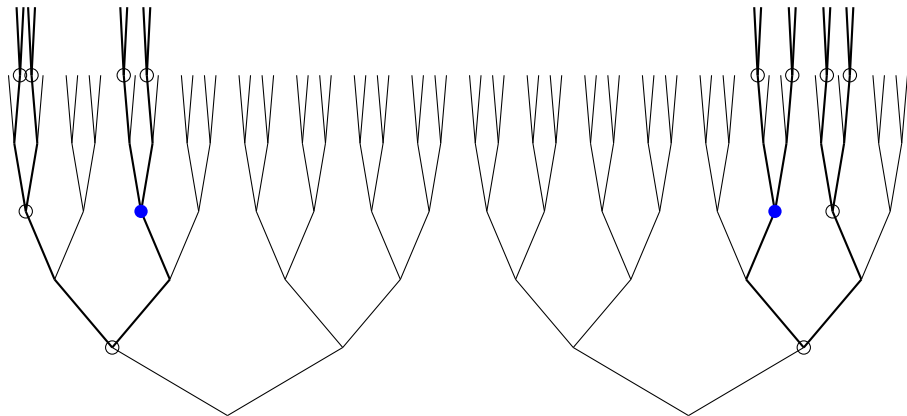
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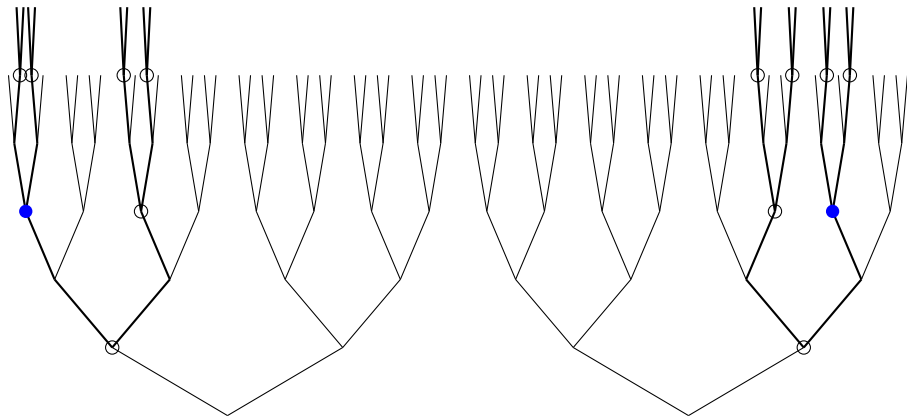


$S_0$

$S_1$



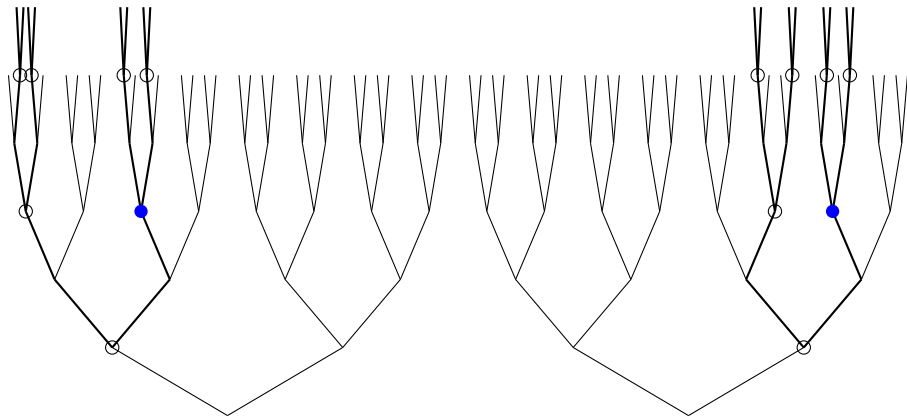
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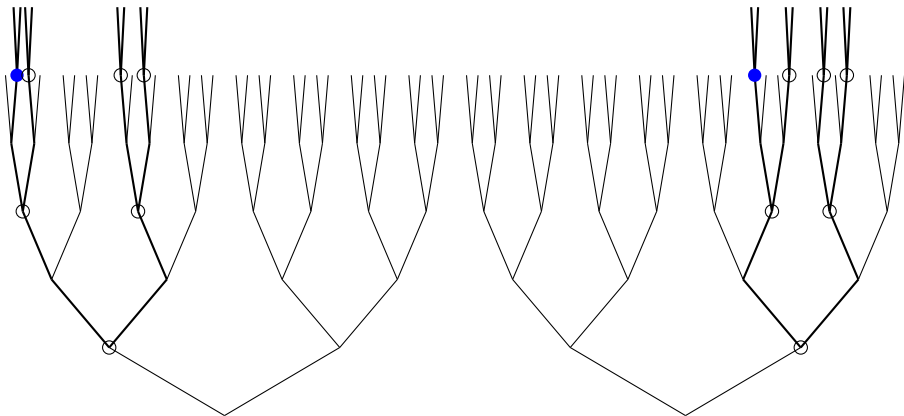
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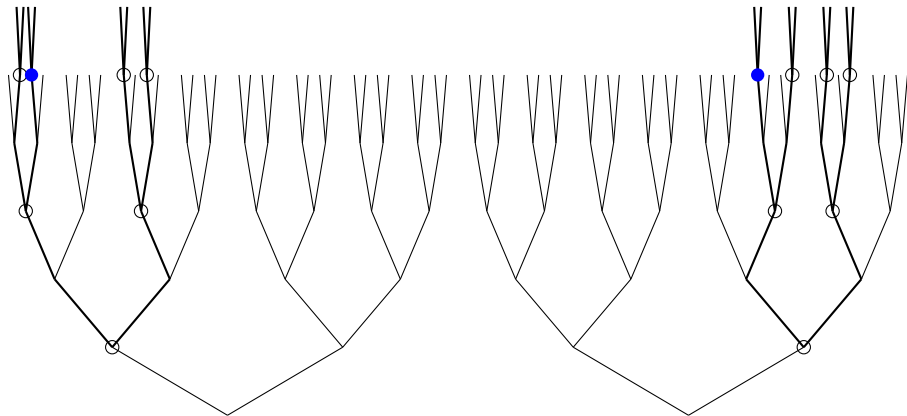
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# HL gives Strong Subtrees with 1 color for level products



$S_0$

$S_1$

## Connection with Forcing

Harrington devised a proof of the Halpern-Läuchli Theorem that uses the method of forcing, though without ever moving into a generic extension of the ground model. This will be important later.

The next theorem is proved by induction from the Halpern-Läuchli Theorem for any finite number of trees.

## A Ramsey Theorem for Strong Trees

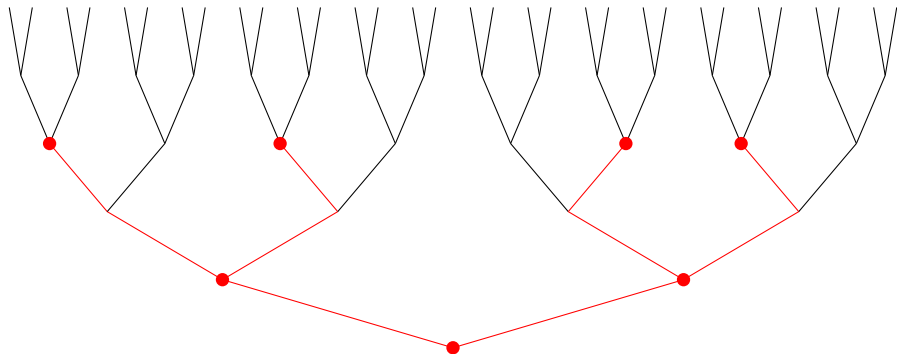
A  $k$ -strong subtree is the truncation of an infinite strong tree to  $k$  many levels.

**Thm.** (Milliken 1979) Let  $k \geq 0$ ,  $l \geq 2$ , and a coloring of all  $k$ -strong subtrees of  $2^{<\omega}$  into  $l$  colors. Then there is an infinite strong subtree  $S \subseteq 2^{<\omega}$  such that all copies of  $2^{\leq k}$  in  $S$  have the same color.

Milliken's theorem for 2-strong trees directly implies the Halpern-Läuchli Theorem.

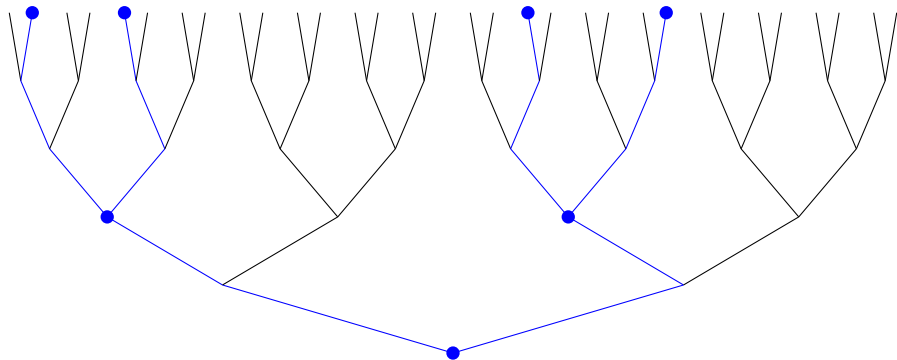
# Milliken's Theorem for 3-Strong Trees

takes a coloring all subtrees of  $2^{<\omega}$  like this:



# Milliken's Theorem for 3-Strong Trees

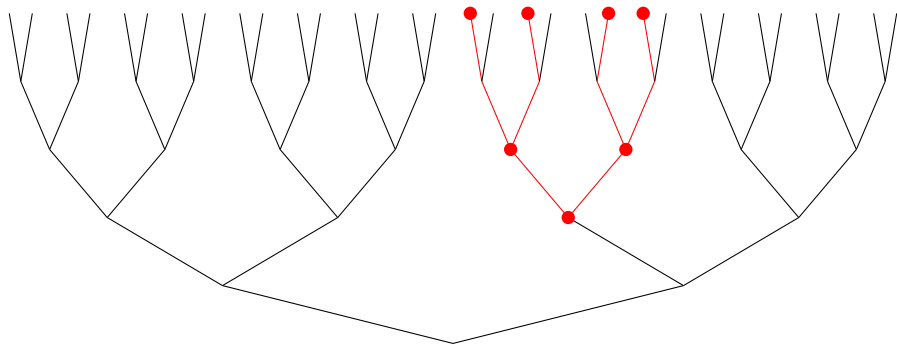
and this:





# Milliken's Theorem for 3-Strong Trees

and this



and finds an infinite strong subtree in which all 3-strong subtrees have the same color.

Applications of this will be seen shortly.

# Extensions of Ramsey's Theorem to Structures

What happens if we try to extend Ramsey's Theorem to infinite structures  $\mathcal{S}$ , where the subsets allowed must have induced structure isomorphic to  $\mathcal{S}$ ?

## Example: The Rationals.

Any finite coloring of the singletons in  $\mathbb{Q}$  is monochromatic on a subset isomorphic to  $\mathbb{Q}$ . However,

**Theorem.** (Sierpinski) There is a coloring of pairs of rationals into two colors such that any subset  $\mathbb{Q}' \subseteq \mathbb{Q}$ , which is again a dense linear order without endpoints, takes both colors on its pairsets.

Decades later, Milliken's Theorem was seen to be the structural heart of this phenomenon.

## Milliken on Ramsey Theory of the Rationals

The rationals can be coded as the nodes in  $2^{<\omega}$ . Applying Milliken's Theorem one finds:

**Fact.** Given any  $n \geq 2$ , there is a number  $T(n, \mathbb{Q}) \geq 2$  such that any coloring of  $[\mathbb{Q}]^n$  into finitely many colors can be reduced to no more than  $T(n, \mathbb{Q})$  colors on a substructure  $\mathbb{Q}'$  isomorphic to  $\mathbb{Q}$ .

With more work, Devlin (building on Laver's work) found the exact numbers: these are tangent numbers! These numbers  $T(n, \mathbb{Q})$  are called **big Ramsey degrees**. They are deduced from the number of **types** of trees that can code an  $n$ -tuple of rationals in  $2^{<\omega}$ .

# Ramsey Theory of the Rado Graph

The **Rado graph**  $\mathcal{R}$  is the universal homogeneous graph on countably many vertices. It is  $\mathcal{R}$  the Fraïssé limit of the class of finite graphs.

**Fact.** (Folklore) Given a coloring of the vertices of the Rado graph into finitely many colors, there is a subgraph which is again Rado in which all vertices have the same color.

The **big Ramsey degree** for vertex colorings in the Rado graph is 1.

# Colorings of Finite Graphs

Example: Ordered graph  $A$  embeds into ordered graph  $B$ .



Figure: Ordered Graph A

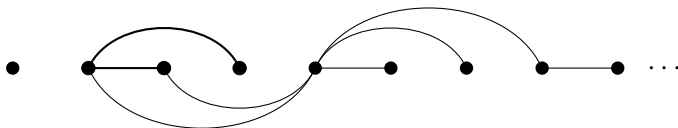
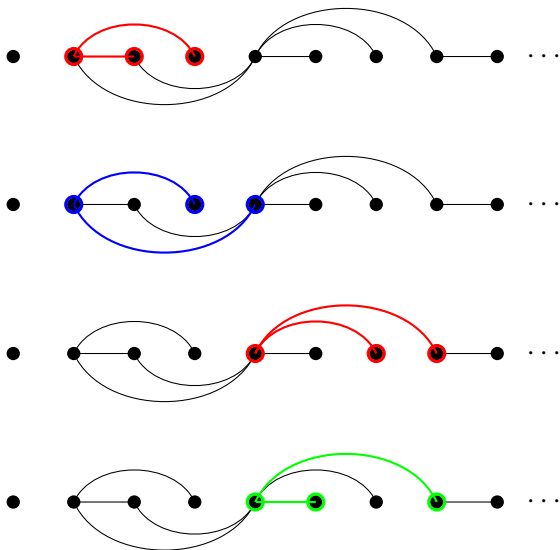


Figure: Ordered Graph B

# Some copies of A in B



# Ramsey Theory of the Rado Graph

- Edges have big Ramsey degree 2. (Pouzet/Sauer 1996).
- All finite graphs have finite big Ramsey degree. (Sauer 2006) In this paper is also the set-up for
- Actual degrees were found structurally in (LSV 2006) and computed in (J. Larson 2008).

How was Milliken's Theorem used?

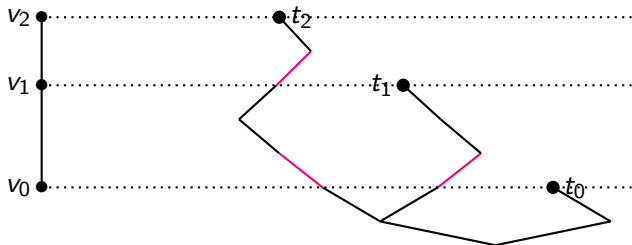
## Nodes in Trees can Code Graphs

Let  $A$  be a graph. Enumerate the vertices of  $A$  as  $\langle v_n : n < N \rangle$ .

A set of nodes  $\{t_n : n < N\}$  in  $2^{<\omega}$  codes  $A$  if and only if for each pair  $m < n < N$ ,

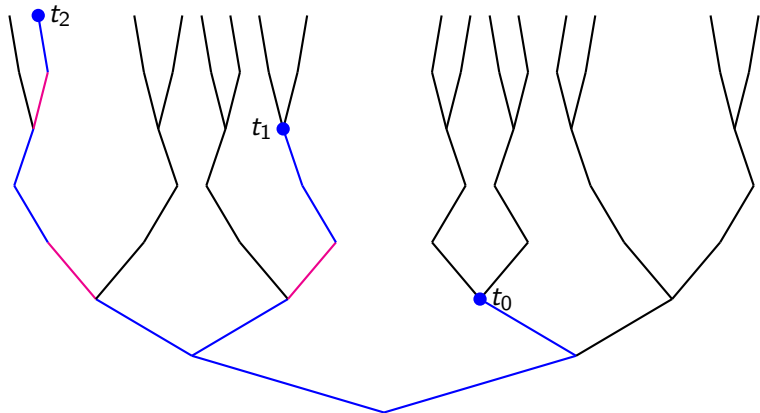
$$v_n E v_m \Leftrightarrow t_n(|t_m|) = 1.$$

The number  $t_n(|t_m|)$  is called the **passing number** of  $t_n$  at  $t_m$ .

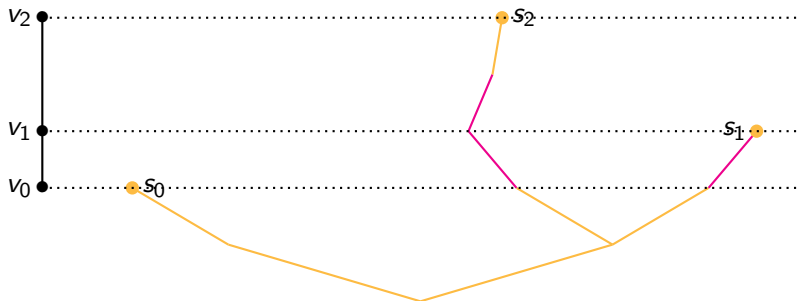




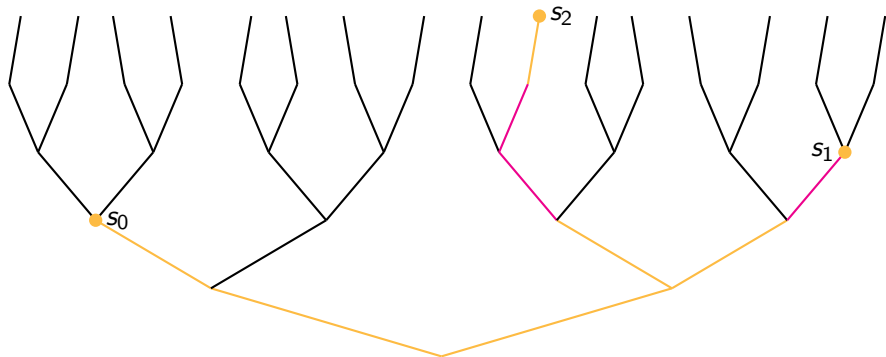
# A Strong Tree Envelope



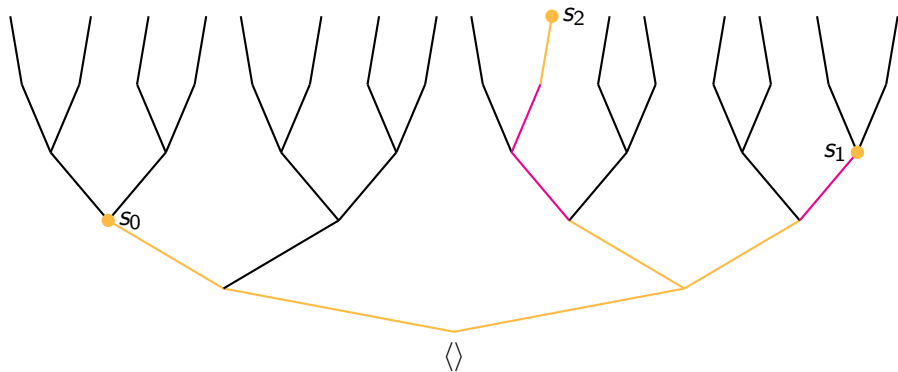
## A Different Antichain Coding a Path of Length 2



# A Strong Tree Envelope



## A different strong tree envelope



# Outline of Sauer's Proof: $\mathcal{R}$ has finite big Ramsey degrees

- 1 The Rado graph is bi-embeddable with the graph coded by all nodes in the tree  $2^{<\omega}$ .
- 2 Each finite graph can be coded by finitely many strong similarity types of (diagonal) antichains.
- 3 Each strongly diagonal antichain can be enveloped into finitely many strong trees.
- 4 Apply Milliken's Theorem finitely many times to obtain one color for each (strong similarity) type.
- 5 Choose a strongly diagonal antichain coding the Rado graph.
- 6 Show that each type persists in each subgraph which is random to obtain exact numbers.

## Structures known to have big Ramsey degrees

- the natural numbers (Ramsey 1929) (all big Ramsey degrees are 1)
- the rationals (Galvin, Laver, Devlin 1979)
- the Rado graph and similar binary relational structures (Sauer 2006)
- the countable ultrametric Urysohn space (Nguyen Van Thé 2008)
- the dense local order, circular tournament,  $\mathbb{Q}_n$  (Laflamme, NVT, Sauer 2010).

The crux of all but two of these proofs is Milliken's Theorem (or variant).

(The Urysohn space result uses Ramsey's Theorem.)

## Missing Piece: Forbidden Configurations

No Fraïssé structure with forbidden configurations had a complete analysis of its Big Ramsey Degrees.

The Problem: Lack of tools for representing such Fraïssé structures and lack of a viable Ramsey theory for such (non-existent) representations.

This problem is addressed starting with my submitted paper, *The Ramsey theory of the universal homogeneous triangle-free graph*, 48 pp, and work-in-progress extending it to all Henson graphs.

The methods developed therein are flexible and should apply, after modifications, to a large collection of homogeneous structures with forbidden configurations.

# Why study Ramsey Theory of Homogeneous Structures?

- Natural extension of structural Ramsey theory on finite structures, and is in line with Ramsey's original theorem.
- Connections with topological dynamics - universal completion flows.
- Possible connections with model theory.



# The Triangle-free Henson Graph $\mathcal{H}_3$ : History of Results

The **universal homogeneous triangle-free graph**  $\mathcal{H}_3$  is the Fraïssé limit of the class of finite triangle-free graphs.

- Henson constructed  $\mathcal{H}_3$  and proved it is weakly indivisible in 1971.
- The Fraïssé class of finite ordered triangle-free graphs has the Ramsey property. (Nešetřil-Rödl 1973)
- $\mathcal{H}_3$  is indivisible: Vertex colorings of  $\mathcal{H}_3$  have big Ramsey degree 1. (Komjáth/Rödl 1986)
- $\mathcal{H}_3$  has big Ramsey degree 2 for edges. (Sauer 1998)

There progress halted due to lack of broadscale techniques.

## Main Theorem: $\mathcal{H}_3$ has Finite Big Ramsey Degrees

**Theorem.** (D.) For each finite triangle-free graph  $A$ , there is a positive integer  $T_{\mathcal{K}_3}(A)$  such that for any coloring of all copies of  $A$  in  $\mathcal{H}_3$  into finitely many colors, there is a subgraph  $\mathcal{H} \leq \mathcal{H}_3$ , again universal triangle-free, such that all copies of  $A$  in  $\mathcal{H}$  take no more than  $T_{\mathcal{K}_3}(A)$  colors.

Thanks to the following:

2011 Laver outlined Harrington's 'forcing proof' of Halpern-Läuchli for me.  
2012 and 2013 Todorćević and Sauer both mention the lack of an appropriate Milliken Theorem as the main obstacle to the solution.

## Structure of Proof: Three Main Parts

- I Develop new notion of **strong coding tree** to represent  $\mathcal{H}_3$ .
- II Prove a Ramsey Theorem for **strictly similar** finite antichains.  

The proof uses ideas from Harrington's 'forcing proof' of the Halpern-Läuchli Theorem, and obtains a Milliken-style theorem.
- III Apply Ramsey Theorem for strictly similar antichains finitely many times. Then take an antichain of coding nodes coding  $\mathcal{H}_3$ .

## Part I: Strong Coding Trees

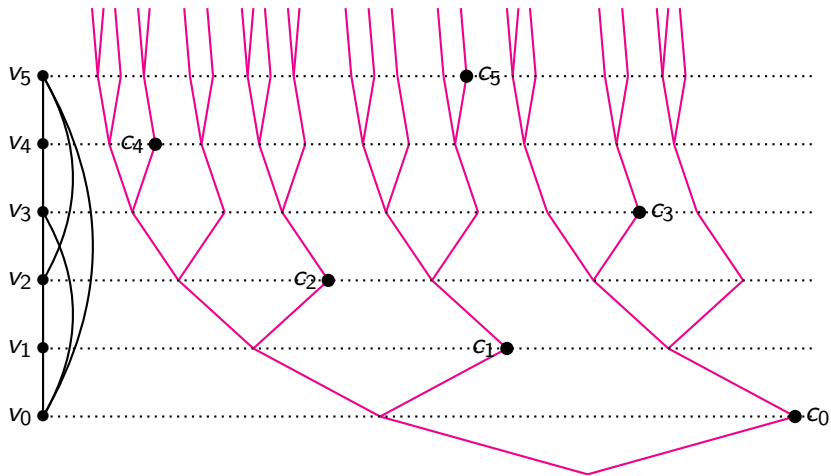
Idea: Want correct analogue of strong trees for setting of  $\mathcal{H}_3$ .

Problem: How to make sure triangles are never encoded but branching is as thick as possible?

## First Approach: Strong Triangle-Free Trees

- Use a **unary predicate** for distinguishing certain nodes to code vertices of a given graph (called **coding nodes**).
- Make a **Branching Criterion** so that a node  $s$  splits iff all its extensions will never code a triangle with coding nodes at or below the level of  $s$ .

# Strong triangle-free tree $\mathbb{S}$

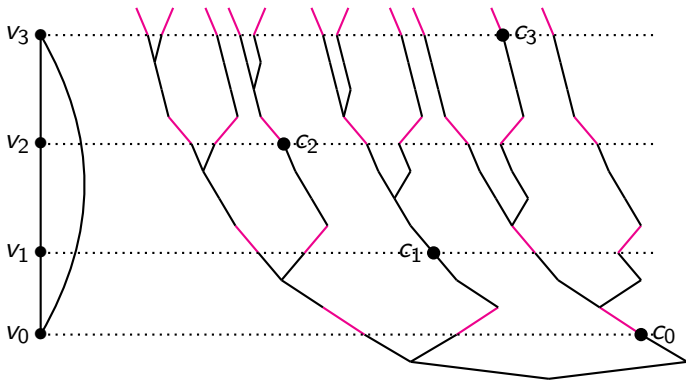


## Almost sufficient

One can develop almost all the Ramsey theory one needs on strong triangle-free trees

**except** for vertex colorings: there is a bad coloring of coding nodes.

## Refined Approach: Strong coding tree $\mathbb{T}$



Skew the levels of interest.



## The Space of Strong Coding Trees: $\mathcal{T}_3$

$\mathcal{T}_3$  is the collection of all subtrees of  $\mathbb{T}$  which are strongly similar to  $\mathbb{T}$ .

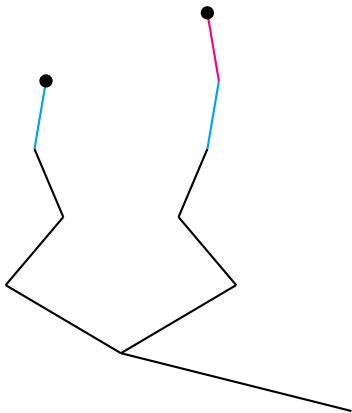
Extension Criterion: A finite subtree  $A$  of a strong coding tree  $T \in \mathcal{T}_3$  can be extended to a strong coding subtree of  $T$  whenever  $A$  is strongly similar to an initial segment of  $\mathbb{T}$  and **all entanglements of  $A$  are witnessed** - no types are lost.

The criteria guaranteeing this are

- 1 **Pre-Triangle Criterion**: All new sets of parallel 1's in  $A$  are witnessed by a coding node in  $A$  'nearby'.
- 2  **$A$  is free in  $T$** :  $A$  has no pre-determined new parallel 1's in  $T$ .

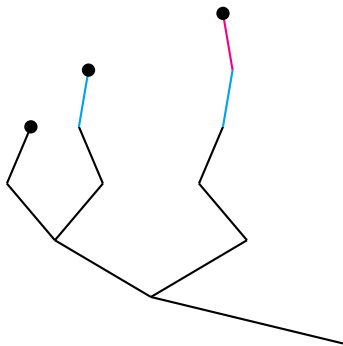
## A subtree of $\mathbb{T}$ in which Pre-Triangle Criterion fails

It has parallel 1's not witnessed by a coding node (PTC fails).



## A subtree of $\mathbb{T}$ in which PTC holds

Its parallel 1's are **witnessed** by a coding node.



This gives the basic idea of PTC, though more subtleties are involved.

## Part II: A Ramsey Theorem for Strictly Similar Finite Antichains.

Idea: Strict similarity takes into account tree isomorphism and placements of coding nodes and new sets of parallel 1's.

It persists upon taking subtrees in  $\mathcal{T}_3$ .

## Ramsey Theorem for Strong Coding Trees

**Theorem.** (D.) Let  $A$  be a finite subtree of a strong coding tree  $T$ , and let  $c$  be a coloring of all strictly similar copies of  $A$  in  $T$ .

Then there is a strong coding tree  $S \leq T$  in which all strictly similar copies of  $A$  in  $S$  have the same color.

This is an analogue of Milliken's Theorem for strong coding trees.

**Strict similarity** is a strong version of isomorphism, and forms an equivalence relation.

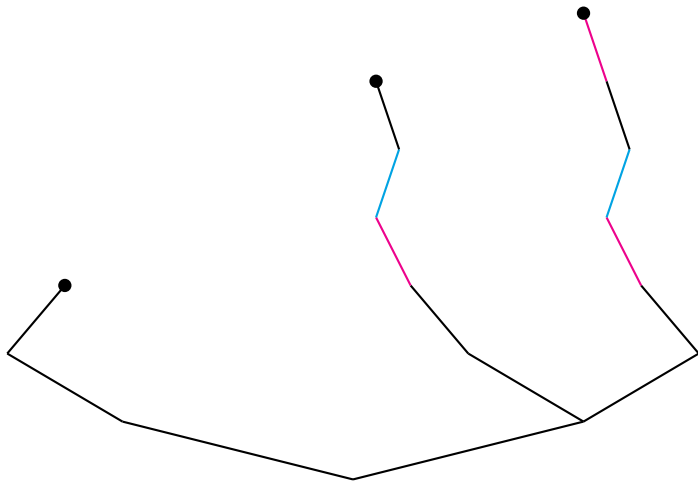
## Some Examples of Strict Similarity Types

Let  $G$  be the graph with three vertices and no edges.

We show some distinct strict similarity types of trees coding  $G$ .

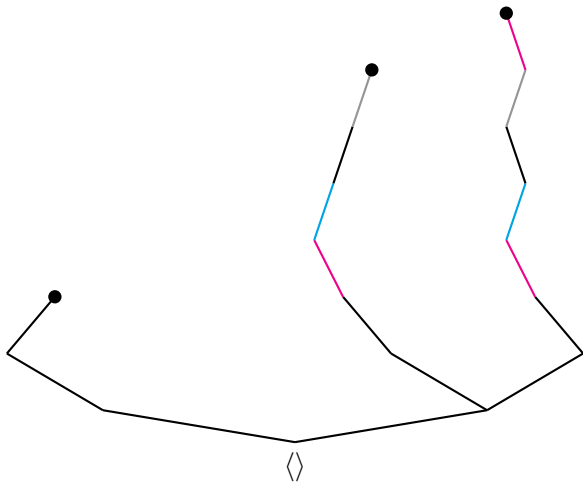
## $G$ a graph with three vertices and no edges

A tree  $A$  coding  $G$  - not P1C but still a valid strict similarity type



# $G$ a graph with three vertices and no edges

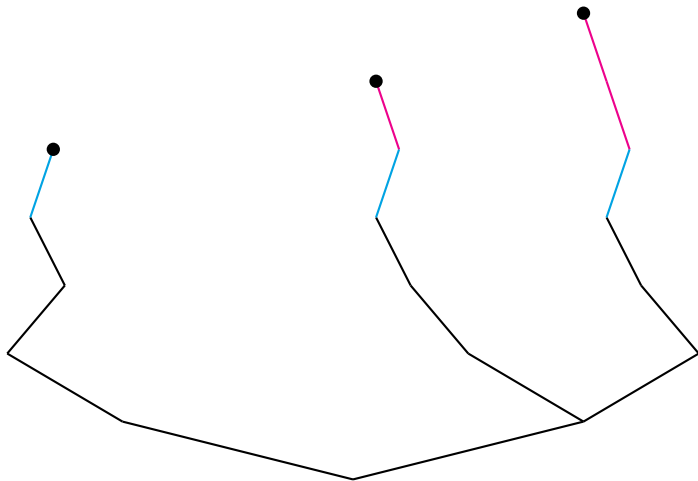
$B$  codes  $G$  and is strictly similar to  $A$ .





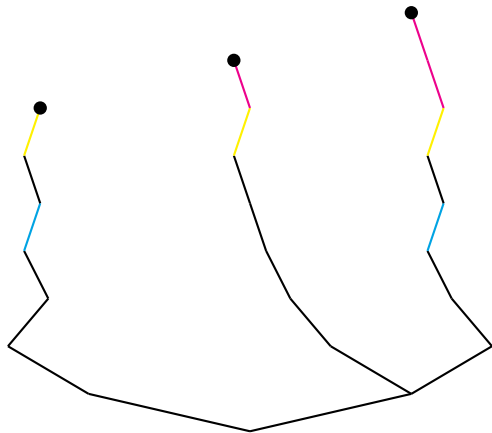
# The tree $C$ codes $G$

$C$  is not strictly similar to  $A$ .

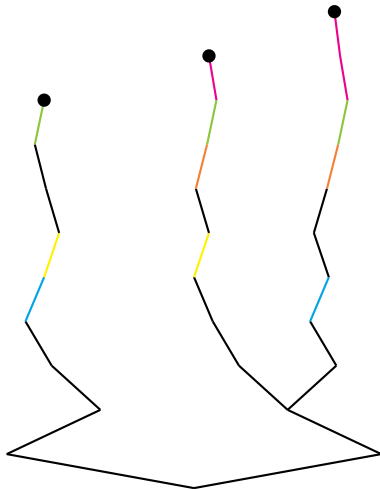


# The tree $D$ codes $G$

$D$  is not strictly similar to either  $A$  or  $C$ .

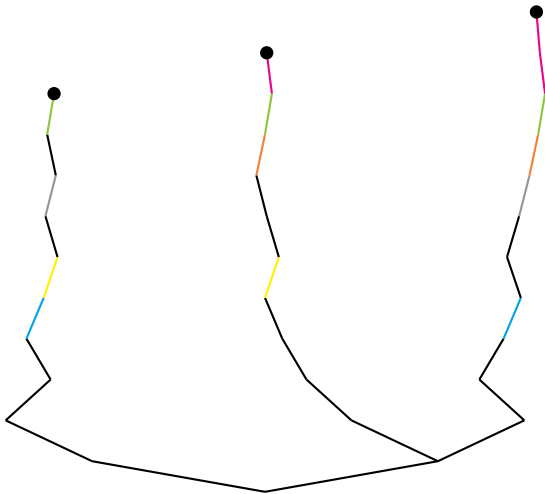


The tree  $E$  codes  $G$  and is not strictly similar to  $A - D$



$E$  is incremental. More on that later.

The tree  $F$  codes  $G$  and is strictly similar to  $E$



$F$  is also incremental.

Part III: Apply the Ramsey Theorem to Strictly Similarity Types  
of Antichains to obtain the Main Theorem.

## Bounds for $T_{\mathcal{K}_3}(G)$

- 1 Let  $G$  be a finite triangle-free graph, and let  $f$  color the copies of  $G$  in  $\mathcal{H}_3$  into finitely many colors.
- 2 Define  $f'$  on antichains in  $\mathbb{T}$ : For an antichain  $A$  of coding nodes in  $\mathbb{T}$  coding a copy,  $G_A$ , of  $G$ , define  $f'(A) = f(G_A)$ .
- 3 List the strict similarity types of antichains of coding nodes in  $\mathbb{T}$  coding  $G$ . There are finitely many.
- 4 Apply the Ramsey Theorem from Part II, once for each strict similarity type, to obtain a strong coding tree  $S \leq \mathbb{T}$  in which  $f'$  has one color per type.
- 5 Take an antichain of coding nodes,  $\mathbb{A}$  in  $S$ , which codes  $\mathcal{H}_3$ . Let  $\mathcal{H}'$  be the subgraph of  $\mathcal{H}_3$  coded by  $\mathbb{A}$ .
- 6 Then  $f$  has no more colors on the copies of  $G$  in  $\mathcal{H}'$  than the number of (incremental) strict similarity types of antichains coding  $G$ .

## Reducing the Upper Bounds

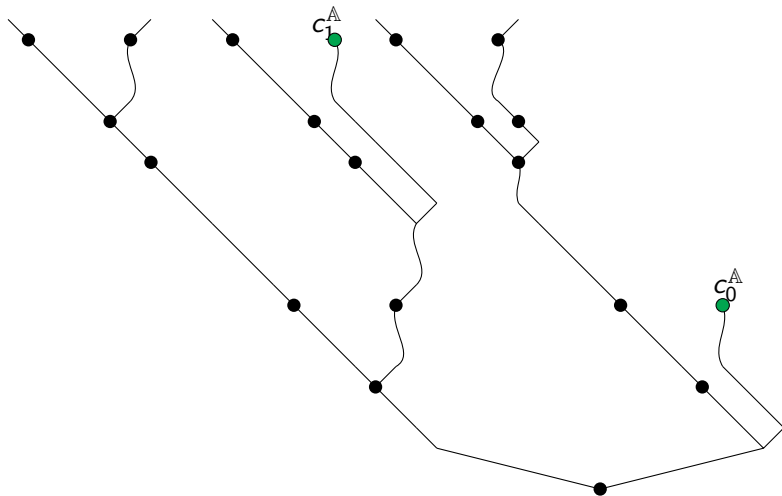
A strong tree  $U$  with coding nodes is **incremental** if whenever a new set of parallel 1's appears in  $U$ , all of its subsets appear as parallel 1's at a lower level.

The trees  $A$ ,  $B$ ,  $E$ , and  $F$  are incremental.

The trees  $C$  and  $D$  are not incremental.

We can take  $S$  in the previous slide to be an incremental strong coding tree.

# An antichain $\mathbb{A}$ of coding nodes of $S$ coding $\mathcal{H}_3$



The tree minus the antichain of  $c_n^{\mathbb{A}}$ 's is isomorphic to  $\mathbb{T}$ .



## Part II Expanded: Ideas behind the proof of the Ramsey Theorem for Strictly Similar Finite Trees

- (a) Prove new Halpern-Läuchli style Theorems for strong coding trees.
  - Three new forcings are needed, but the proofs take place in ZFC.
- (b) Prove a new Ramsey Theorem for finite trees satisfying the Strict Pre-Triangle Criterion.
  - An analogue of Milliken's Theorem.
- (c) New notion of envelope.
  - Turns an antichain into a tree satisfying Strict Pre-Triangle Criterion.

## (a) Halpern-Läuchli-style Theorem

**Thm.** (D.) Given a strong coding tree  $T$  and

- ①  $B$  a finite, valid strong coding subtree of  $T$ ;
- ②  $A$  a finite subtree of  $B$  with  $\max(A) \subseteq \max(B)$ ; and
- ③  $X$  a level set extending  $A$  into  $T$  with  $A \cup X$  satisfying the PTC and valid in  $T$ .

Color all end-extensions  $Y$  of  $A$  in  $T$  for which  $A \cup Y$  is strictly similar to  $A \cup X$  into finitely many colors.

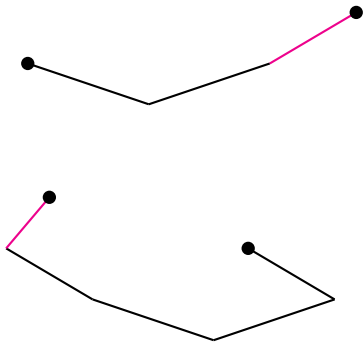
Then there is a strong coding tree  $S \leq T$  end-extending  $B$  such that all level sets  $Y$  in  $S$  with  $A \cup Y$  strictly similar to  $A \cup X$  have the same color.

**Remark.** The proof uses three different forcings and Harrington-style ideas. The forcings are best thought of as conducting unbounded searches for finite objects in ZFC.

Proving the lower bounds in general for big Ramsey degrees of  $\mathcal{H}_3$  is a work in progress.

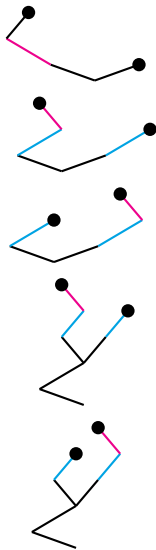
Big Ramsey degrees for edges and non-edges have been computed.

## Edges have big Ramsey degree 2 in $\mathcal{H}_3$



$T_{\mathcal{H}_3}(\text{Edge}) = 2$  was obtained in (Sauer 1998) by different methods.

## Non-edges have 5 Strict Similarity Types (D.)



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