

RETRACT ORTHOGONALITY AND ORTHOGONALITY OF n -ARY OPERATIONS

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CONSTRUCTION METHODS OF ORTHOGONAL OPERATIONS

- [1] **A.S. Bektenov, T Yakubov**, *Systems of orthogonal n -ary operations*, Izvestiya AN MSSR, Seria fiz.-tehn. i mat. nayk, 1974, **3**, 7 – 17.
- [2] **M. Trenkler**, *On orthogonal latin p -dimensional cubes*, Czechoslovak Mathematical Journal, 2005, **55 (130)**, 725 – 728.
- [3] **G. Belyavskaya, G.L. Mullen**, *Orthogonal hypercubes and n -ary operations*, Quasigroups and Related Systems, 2005, **13**, 1, 73 – 86.
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Definition of invertibility

An operation f is called *i-invertible* if for arbitrary elements $a_1, \dots, a_{i-1}, b, a_{i+1}, \dots, a_n$ there exists a unique element x such that

$$f(a_1, \dots, a_{i-1}, x, a_{i+1}, \dots, a_n) = b. \quad (3)$$

If f is *i-invertible* for all $i \in \overline{1, n} := \{1, 2, \dots, n\}$, then it is called an *invertible* or *quasigroup operation*.

- 1 relations between orthogonal n -ary operations and their orthogonal retracts;
- 2 complementing of a k -tuple of orthogonal n -ary operations ($k < n$) to an n -tuple of orthogonal n -ary operations.

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Definition of retract

Let f be an n -ary operation on a set Q .

An $(n - 1)$ -ary operation $f_{i,a}$ is called a *retract* of f by element a , if it is obtained from f by replacing variable x_i with an element $a \in Q$, i.e., if it is defined by

$$f_{i,a}(x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_n) = f(x_1, \dots, x_{i-1}, a, x_{i+1}, \dots, x_n). \quad (4)$$

Definition of retract

Unary operation $f_{(\bar{b}, \{i\})}$ defined by

$$f_{(\bar{b}, \{i\})}(x_i) = f(b_1, \dots, b_{i-1}, x_i, b_{i+1}, \dots, b_n) \quad (5)$$

is unary $(\bar{b}, \{i\})$ -retract of f , $\bar{b} := (b_1, \dots, b_{i-1}, b_{i+1}, \dots, b_n)$.

$$\delta := \{i_1, \dots, i_k\} \subseteq \overline{1, n}, \quad \{j_1, \dots, j_{n-k}\} := \overline{1, n} \setminus \delta, \quad \bar{a} := (a_{j_1}, \dots, a_{j_{n-k}}).$$

An operation $f_{(\bar{a}, \delta)}$ which is defined by

$$f_{(\bar{a}, \delta)}(x_{i_1}, \dots, x_{i_k}) := f(y_1, \dots, y_n), \quad (6)$$

where $y_i := \begin{cases} x_i, & \text{if } i \in \delta, \\ a_i, & \text{if } i \notin \delta \end{cases}$ is called (\bar{a}, δ) -retract or δ -retract of f .

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Definition of retract orthogonality

Definition

Operations $f_{1;(\bar{a}_1,\delta)}$, $f_{2;(\bar{a}_2,\delta)}$, \dots , $f_{k;(\bar{a}_k,\delta)}$ are called *similar δ -retracts* of n -ary operations f_1, \dots, f_k if $\bar{a}_1 = \bar{a}_2 = \dots = \bar{a}_k$.

Definition

If all similar δ -retracts of f_1, \dots, f_k are orthogonal, then the operations f_1, \dots, f_k are called *δ -retractly orthogonal*.

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Problem 1. Relations between orthogonality and retract orthogonality

Theorem 1 (I. Fryz, 2017)

If for some $\delta \in \overline{1, n}$ a tuple of n -ary operations is δ -retractly orthogonal, then the tuple is orthogonal.

Theorem 2 (I. Fryz, 2017)

There exist k -tuples of orthogonal n -ary operations such that for some $\delta \in \overline{1, n}$, $|\delta| = k$, they are not δ -retractly orthogonal.

I.Fryz, *Orthogonality and retract orthogonality of operations* (in print).

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Specification for central quasigroups

If an n -ary quasigroup f is linear on a group $(Q; +)$, then

$$f(x_1, \dots, x_n) = \alpha_1 x_1 + \dots + \alpha_n x_n + a, \quad (7)$$

where $a \in Q$ and $\alpha_1, \dots, \alpha_n$ are automorphisms of $(Q; +)$. If $(Q; +)$ is abelian, then f is called a *central quasigroup* (or a *T-quasigroup*).

Theorem 3 (I. Fryz, 2017)

Let $k \leq n$ and p be a prime number. n -ary central quasigroups f_1, \dots, f_k over field $(\mathbb{Z}_p; +, \cdot)$ are orthogonal if and only if there exists δ such that $|\delta| = k$ and f_1, \dots, f_k are δ -retractly orthogonal.

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Problem 2. Complementing orthogonal operations

Embedding of orthogonal operations

Theorem (G. Belyavskaya, G.Mullen, 2005)

Every k -tuple of orthogonal n -ary operations ($k < n$) can be embedded in an n -tuple of orthogonal n -ary operations.

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This means that for every k -tuple of orthogonal n -ary operations f_1, \dots, f_k , there exist an $(n - k)$ -tuple of orthogonal n -ary operations f_{k+1}, \dots, f_n such that n -tuple f_1, \dots, f_n is orthogonal.

Every complete n -ary operation is complementable to an n -tuple of orthogonal n -ary operations.

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Complementing orthogonal n -ary operations

Algorithm 1.

Let $\delta = \{i_1, \dots, i_k\} \subset \overline{1, n}$ and g_{i_1}, \dots, g_{i_k} are δ -retractly orthogonal Operations $g_{i_{k+1}}, \dots, g_{i_n}$ are constructed by

- 1) choose a partition $\pi := \{\delta, \pi_2, \dots, \pi_q\}$ of $\overline{1, n}$, $f_{i_{k+1}}, \dots, f_{i_n}$ such that for every $r \in \overline{2, q}$ a tuple $\{f_j | j \in \pi_r\}$ is π_r -retractly orthogonal;
- 2) for every $j \in \pi_2$, operation g_j is constructed by

$$g_j(x_1, \dots, x_n) := f_j(t_1, \dots, t_n), \quad (8)$$

where

$$t_s := \begin{cases} g_s(x_1, \dots, x_n), & \text{if } s \in \delta, \\ x_s, & \text{otherwise;} \end{cases}$$

- r) for every $j \in \pi_r$, $r = 3, \dots, q$, operation g_j is constructed by (8), where

$$t_s := \begin{cases} g_s(x_1, \dots, x_n), & \text{if } s \in \delta \cup \pi_2 \cup \dots \cup \pi_{r-1}, \\ x_s, & \text{otherwise.} \end{cases}$$

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Complementing orthogonal n -ary operations

Theorem 4 (I. Fryz, 2017)

A $(n - k)$ -tuple of n -ary operations $g_{i_{k+1}}, \dots, g_{i_n}$ constructed by Algorithm 1 is a complement of a k -tuple of δ -retractly orthogonal n -ary operations g_{i_1}, \dots, g_{i_k} to an n -tuple of orthogonal n -ary operations.

If partition $\pi := \{\delta, \pi_1, \dots, \pi_q\}$ of $\overline{1, n}$, where $\pi_r =: \{i_{k+r}\}$ for all $r \in \overline{1, n - k}$, then constructed complements are trivial. Thus, we have to take i_{k+1}, \dots, i_n -invertible n -ary operations.

Theorem 6 (I. Fryz, 2017)

The number of all complements constructed by Algorithm 1 of a k -tuple of δ -retractly orthogonal n -ary operations ($|\delta| = k$, $k < n$) on Q of order m to an n -tuple of orthogonal n -ary operations is greater than

$$\frac{(m!)^{(n-k)m^{n-1}}}{(n-k)!}.$$

Theorem 5 (I. Fryz, 2017)

Algorithm 2 complements a k -tuple of orthogonal k -ary operations to an n -tuple of orthogonal n -ary operations.

Theorem (I. Fryz, 2017)

The number of all complements constructed by Algorithm 2 of a k -tuple of orthogonal k -ary operations on a set Q of order m to an n -tuple of orthogonal n -ary operations is greater than

$$\frac{(m!)^{km^{n-k}+(n-k)m^{n-1}}}{(n-k)!}.$$

Thank you for your attention